

# $p$ -adic numerical methods in arithmetic geometry

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# Algebraic curves

An *algebraic variety* is the zero locus of finite number of polynomial equations in finite number of variables.

Defined *over* a field  $k$  if the polynomials have coefficients in that field.

A *curve* is an algebraic variety of dimension 1, always smooth and projective, e.g.  $y^2 = x^3 + ax + b$  elliptic curve.

However, sometimes defined by (*singular*) plane model  $f(x, y) = 0$ .

$k$  will always be a finite field  $\mathbb{F}_q$  or the field of rational numbers  $\mathbb{Q}$ .

Recall that  $q = p^n$  and for all such  $q$  there exists a unique finite field  $\mathbb{F}_q$  with  $q$  elements.

# Zeta function

Let  $X$  be an algebraic curve over a finite field  $\mathbb{F}_q$  with  $q = p^n$ .

## Definition

$$Z(X, T) := \exp \left( \sum_{i=1}^{\infty} |X(\mathbb{F}_{q^i})| \frac{T^i}{i} \right)$$

## Example (Projective line)

$$\begin{aligned} Z(\mathbb{P}_{\mathbb{F}_q}^1, T) &= \exp \left( \sum_{i=1}^{\infty} (q^i + 1) \frac{T^i}{i} \right) \\ &= \exp \left( \sum_{i=1}^{\infty} \frac{T^i}{i} \right) \exp \left( \sum_{i=1}^{\infty} \frac{(qT)^i}{i} \right) \\ &= \frac{1}{(1 - T)(1 - qT)} \end{aligned}$$

# Weil conjectures

Theorem (Weil, 1948)

Let  $X$  be a smooth projective curve of genus  $g$  over  $\mathbb{F}_q$ . Then

$$Z(X, T) = \frac{\chi(T)}{(1-T)(1-qT)}$$

with  $\chi(T) \in \mathbb{Z}[T]$  of degree  $2g$ . Moreover,

$$\chi(T) = \prod_{i=1}^{2g} (1 - \alpha_i T),$$

where:

- the  $\alpha_i$  are algebraic integers,
- of absolute value  $\sqrt{q}$ ,
- permuted by  $\alpha \mapsto q/\alpha$ .

## Computing zeta functions

Since the zeta function is given by a *finite* amount of data, one can hope to compute it.

### Problem

Compute  $Z(X, T)$  efficiently.

Bounds on the the degrees of the numerator and denominator of  $Z(X, T)$  are known, so computing  $Z(X, T)$  reduces to computing a finite number of  $X(\mathbb{F}_{q^i})$ .

For a curve of genus  $g$ , have to compute up to  $X(\mathbb{F}_{q^g})$ . Counting *naively* need at least  $q^g$  operations. Too slow for all but the smallest values of  $q$  and  $g$ .

Let us first give some applications.

# Cryptography

Can associate to curve  $X/\mathbb{F}_q$  a finite Abelian group  $J(\mathbb{F}_q)$  called its *Jacobian*. The order of this group is  $\chi(1)$ . The *Discrete Logarithm Problem* (DLP) on  $J(\mathbb{F}_q)$  is:

## Problem

given  $P, Q \in J(\mathbb{F}_q)$  find (the smallest)  $n \in \mathbb{N}$  such that  $nP = Q$ .

This problem is used in cryptography in *Diffie Helmann* key exchange. When the order of  $J(\mathbb{F}_q)$  only has small prime factors the DLP is easy.

So we need to compute  $\chi(1)$  and we can do this by computing  $Z(X, T)$ .

## Sato-Tate distributions

Let  $X$  be a smooth projective curve of genus  $g$  defined over  $\mathbb{Q}$ .

For every prime  $p$  let  $X_p/\mathbb{F}_p$  denote the reduction of  $X$  modulo  $p$ . Again, for all but a finite number of  $p$ :

$$Z(X_p, T) = \frac{\chi_p(T)}{(1-T)(1-qT)}$$

for some polynomial  $\chi_p(T) \in \mathbb{Z}[T]$  of degree  $2g$ .

### Problem

*How is the polynomial  $\chi_p(T/\sqrt{p})$  distributed when  $p$  varies?*

*Conjectural answer:* as the (reverse) characteristic polynomial of a random conjugacy class of a certain compact group. So far only known for  $g = 1$ .

Andrew Sutherland (with coauthors) computed  $\chi_p(T)$  for  $X$  with  $g = 2$  and found all predicted distributions!

# $p$ -adic numbers

Let:

- $p$  a prime,
- $\text{ord}_p$  the  $p$ -adic valuation on  $\mathbb{Z}$  (number of factors  $p$ ),
- $\|\cdot\|_p = p^{-\text{ord}_p(\cdot)}$  associated metric on  $\mathbb{Z}$ .

## Definition ( $p$ -adic integers)

$\mathbb{Z}_p$  is the completion of  $\mathbb{Z}$  w.r.t.  $\|\cdot\|_p$ .

Elements  $a_0 + a_1p + a_2p^2 + \dots$ , with  $a_i \in \{0, 1, \dots, p-1\}$ .

$\mathbb{Z}_p$  has a unique maximal ideal  $(p)$  and  $\mathbb{Z}_p/p\mathbb{Z}_p \cong \mathbb{F}_p$ . The field of fractions  $\mathbb{Q}_p$  has characteristic 0.

Similarly, can define  $\mathbb{Z}_q$  such that  $\mathbb{Z}_q/p\mathbb{Z}_q \cong \mathbb{F}_q$ , field of fractions  $\mathbb{Q}_q$ .



# $p$ -adic cohomology

Let  $X/\mathbb{F}_q$  be a smooth projective curve.

Can define  $p$ -adic (rigid/Monsky Washnitzer) cohomology space  $H_{\text{rig}}^1(X)$ :

- finite dimensional  $\mathbb{Q}_q$  vector space,
- with action  $F_q^*$  of the  $q$ -th power map  $F_q$ ,

such that:

$$\chi(T) = \det(1 - F_q^* T | H_{\text{rig}}^1(X)).$$

**Idea:** first lift  $X$  to characteristic 0, then take (overconvergent) de Rham cohomology.

## Smooth affine case

Let:

- $X = \text{Spec}(\bar{A})$  with  $\bar{A}$  a smooth  $\mathbb{F}_q$  algebra of finite type.
- smooth lift  $A = \mathbb{Z}_q[x_1, \dots, x_l]/(f_1, \dots, f_m)$  such that  $A \otimes_{\mathbb{Z}_q} \mathbb{F}_q = \bar{A}$ .
- weak completion  $A^\dagger = \mathbb{Z}_q\langle x_1, \dots, x_l \rangle^\dagger / (f_1, \dots, f_m)$ , where

$$\mathbb{Z}_q\langle x_1, \dots, x_l \rangle^\dagger := \left\{ \sum_I a_I x^I : a_I \in \mathbb{Z}_q, \exists \rho > 1 \text{ s.t. } \lim_{|I| \rightarrow \infty} |a_I| \rho^{|I|} = 0 \right\}.$$

- 1-forms  $\Omega_{A^\dagger}^1 = (A^\dagger dx_1 \oplus \dots \oplus A^\dagger dx_l) / (A^\dagger df_1 + \dots + A^\dagger df_m)$ .
- $d : A^\dagger \mapsto \Omega_{A^\dagger}^1$ .

Then:  $H_{\text{rig}}^1(X) = \text{coker}(d) \otimes \mathbb{Q}_q$ .

# Hyperelliptic curves

Let:

- $\mathbb{F}_q$  a finite field of odd characteristic,
- $Q \in \mathbb{F}_q[x]$  monic of degree  $2g + 1$  without repeated roots,
- $X$  the smooth projective curve defined by  $y^2 = Q(x)$ .

$X$  is a hyperelliptic curve of genus  $g$  (with a rational Weierstrass point).

For characteristic 2 or curves without a rational Weierstrass point this has to be modified slightly.

Kedlaya (2001) proposed to compute  $Z(X, T)$  using  $p$ -adic cohomology.

# Kedlaya's algorithm

Sketch:

- $X : y^2 = Q(x)$  hyperelliptic of genus  $g$  over  $\mathbb{F}_q$  with  $q = p^n$  odd.
- $U$  open in  $X$  defined by  $y \notin \{0, \infty\}$ .
- $Q \in \mathbb{Z}_q[x]$  a monic lift of  $Q$ .
- $A^\dagger = \mathbb{Z}_q\langle x, y, 1/y \rangle^\dagger / (y^2 - Q)$ .
- basis for  $H_{rig}^1(X) \subset H_{rig}^1(U)$  given by  $[\frac{dx}{y}, \dots, x^{2g-1} \frac{dx}{y}]$ .
- lift Frobenius to  $A^\dagger$ :  $F_p(x) := x^p$ , find  $F_p(y) \equiv y^p \pmod{p}$  (Hensel).
- Apply  $F_p$  to basis for  $H_{rig}^1(X)$  and reduce to find matrix  $F_p$ .
- Compute  $\chi(T) = \det(1 - TF_p^n | H_{rig}^1(X))$ .
- $Z(X, T) = \chi(T) / ((1 - T)(1 - qT))$ .

# Complexity

$X/\mathbb{F}_q$  hyperelliptic  $q = p^n$  genus  $g$ .

Theorem (Kedlaya, 2001)

$Z(X, T)$  can be computed in time  $O((pg^4n^3)^{1+\epsilon})$ .

Input size about  $\log(p)gn$ , so (only) polynomial time for fixed  $p$ .

All  $p$ -adic algorithms suffer from this, but the dependence on  $p$  can be improved to  $O(p^{1/2+\epsilon})$  and *average polynomial time* (Harvey).

## General curves

Let  $X/\mathbb{F}_q$  with  $q = p^n$  be the smooth projective curve birational to a (singular) plane curve

$$f(x, y) = 0$$

with  $f \in \mathbb{F}_q[x, y]$  irreducible and monic in  $y$  of degree  $d_x, d_y$  in  $y, x$ .

Theorem (Tuitman, 2014)

*Suppose that we know a 'good' lift  $F \in \mathbb{Z}_q[x, y]$  of  $f$  to characteristic zero (technical). Then the zeta function of  $X$  can be computed in time:*

$$O((pd_y^6 d_x^4 n^3)^{1+\epsilon})$$

Input size about  $\log(p)d_x d_y n$ , so again (only) polynomial time for fixed  $p$ .

Good lift condition rather technical. However often known or easy to construct (e.g. when  $g \leq 5$ , joint work with Castryck).

# Some examples

## Example: the modular curve $X_1(17)$

```
> P<x>:=PolynomialRing(RationalField());
> R<y>:=PolynomialRing(P);
> f:=y^4 + (x^3 + x^2 - x + 2)*y^3 + (x^3 - 3*x + 1)*y^2 - (x^4 + 2*x)*y + x^3 + x^2;
> p:=101;
> ZetaFunction(f,p);
(10510100501*T^10 + 6035503258*T^9 + 1900905345*T^8 + 396288448*T^7 + 60231754*T^6 + 6865620*T^5 +
596354*T^4 + 38848*T^3 + 1845*T^2 + 58*T + 1)/(101*T^2 - 102*T + 1)
```

## Example: a generic genus 5 curve

```
> C:=RandomGenus5CurveNonTrigonal(FiniteField(37));
> C;
Curve over GF(37) defined by
19*$$.1^2 + 18*$$.1*$.2 + 31*$$.2^2 + $.1*$.3 + 19*$$.2*$.3 + 25*$$.3^2 + 8*$$.1*$.4 + 17*$$.2*$.4 + 29*$$.3*$.4 +
19*$$.4^2 + 18*$$.1*$.5 + 27*$$.2*$.5 + 26*$$.3*$.5 + 14*$$.4*$.5 + 32*$$.5^2,
12*$$.1^2 + 31*$$.1*$.2 + 18*$$.2^2 + 11*$$.1*$.3 + 24*$$.2*$.3 + 21*$$.3^2 + 12*$$.1*$.4 + 4*$$.2*$.4 + 21*$$.3*$.4 +
22*$$.4^2 + 4*$$.1*$.5 + 31*$$.2*$.5 + 23*$$.3*$.5 + 20*$$.4*$.5 + 35*$$.5^2,
21*$$.1^2 + 35*$$.1*$.2 + 17*$$.2^2 + 8*$$.1*$.3 + 12*$$.2*$.3 + 32*$$.3^2 + 34*$$.1*$.4 + 22*$$.2*$.4 + 24*$$.3*$.4 +
18*$$.4^2 + 19*$$.1*$.5 + 10*$$.2*$.5 + 19*$$.3*$.5 + 10*$$.4*$.5
> ZetaFunction(C);
(69343957*T^10 - 5622483*T^9 + 1418284*T^8 + 217671*T^7 - 2997*T^6 + 6604*T^5 - 81*T^4 + 159*T^3 + 28*T^2
- 3*T + 1)/(37*T^2 - 38*T + 1)
```

Fields very small in these examples, but same thing works over  $GF(p)$  with  $p \sim 2^{15}$  and  $GF(q)$  with  $q = p^n$  much larger still!

# Precision

Can only compute with  $p$ -adic numbers to finite precision  $N$  (i.e. mod  $p^N$ ).

Recall that

$$\chi(T) = \chi_{2g} T^{2g} + \dots + \chi_1 T + 1 = \prod_{i=1}^{2g} (1 - \alpha_i T)$$

with  $|\alpha_i| = \sqrt{q}$  and  $\alpha \mapsto q/\alpha$  permuting the  $\alpha_i$ . So  $|\chi_i| \leq \binom{2g}{g} q^{g/2}$  and  $\chi_{2g-i} = q^{g-i} \chi_i$ . Sufficient to compute  $\chi$  to precision  $N$  such that

$$p^N > 2 \binom{2g}{g} q^{g/2}.$$

Have to keep track of  $p$ -adic precision throughout the algorithm. This involves very interesting (and rather technical) mathematics.



## Rational points

$X/\mathbb{Q}$  a smooth projective curve of genus  $g > 1$ .

Given by (singular) plane model  $f(x, y) = 0$ .

Theorem (Faltings, 1983)

*The set  $X(\mathbb{Q})$  of rational points on  $X$  is finite.*

Usually points are easily found by a search (if they exist).

Example ( $g = 4$ )

$$f(x, y) = y^3 - (x^5 - 2x^4 - 2x^3 - 2x^2 - 3x)$$

$$X(\mathbb{Q}) \supset \{(1, -2), (0, 0), (-1, 0), (3, 0), \infty\}$$

Problem

*How to prove that these are all points?*

# Chabauty's theorem

$J$  will denote the Jacobian variety of  $X$ , i.e. divisors of degree 0 modulo divisors of functions. Note that  $J$  is naturally an abelian variety.

## Theorem (Mordell-Weil)

$J(\mathbb{Q})$  is a finitely generated abelian group.

Given a point  $b \in X(\mathbb{Q})$ , we get an embedding  $X(\mathbb{Q}) \rightarrow J(\mathbb{Q})$ :

$$P \mapsto (P) - (b)$$

## Theorem (Chabauty, 1941)

Let  $r$  be the rank of  $J(\mathbb{Q})$ . If  $r < g$  then  $X(\mathbb{Q})$  is finite.

Coleman: can make this effective using  $p$ -adic line integrals.

# Coleman integrals

Let:

- $p$  a prime at which  $X$  has good reduction,
- $P, Q \in X(\mathbb{Q}_p)$ ,
- $\omega$  a 1-form on  $X_{\mathbb{Q}_p}$  (more generally on some wide open of a rigid analytic space).

In the 80's Coleman defined path independent line integrals

$$\int_P^Q \omega$$

which can be extended to integrate over  $D \in J(\mathbb{Q}_p)$ , where  $J$  is the Jacobian of  $X$  (above:  $D = (Q) - (P)$ ).

# Properties

The Coleman integral has the following properties:

- 1 Linearity:  $\int_P^Q (a\omega_1 + b\omega_2) = a \int_P^Q \omega_1 + b \int_P^Q \omega_2$ .
- 2 Additivity in endpoints:  $\int_P^Q \omega = \int_P^R \omega + \int_R^Q \omega$ .
- 3 Change of variables: If  $V' \subset X'$  is a wide open subspace of a rigid analytic space  $X'$  and  $\phi : V \rightarrow V'$  a rigid analytic map then  $\int_P^Q \phi^* \omega = \int_{\phi(P)}^{\phi(Q)} \omega$ .
- 4 Fundamental theorem of calculus:  $\int_P^Q df = f(Q) - f(P)$  for  $f$  a rigid analytic function on  $V$ .

A residue disk on  $X_{\mathbb{Q}_p}$  is the inverse image under reduction mod  $p$  of a single point. Coleman integrals within a single residue disk are called tiny.

## Tiny integrals

Let:  $P, Q \in X(\mathbb{Q}_p)$  points in the same residue disk,  $\omega \in H^0(X_{\mathbb{Q}_p}, \Omega^1)$ .

Then  $\int_P^Q \omega$  can be computed by expanding  $\omega$  in a local coordinate  $t$  on the disk:

$$\omega = \sum_{i \geq 0} c_i t^i dt$$

and integrating as usual

$$\int_{t(P)}^{t(Q)} \sum_{i \geq 0} c_i t^i dt = \sum_{i \geq 0} \frac{c_i}{i+1} (t(Q)^{i+1} - t(P)^{i+1}).$$

When  $P$  and  $Q$  not in the same residue disk, does not work: series do not converge.

Analytic continuation fails over  $\mathbb{Q}_p$ . Coleman: use Frobenius action on  $p$ -adic cohomology.

## $p$ -adic cohomology

Let  $U \subset X$  be an open such that  $X - U$  is smooth over  $\mathbb{Z}_p$  and  $\omega_1, \dots, \omega_{2g} \in \Omega^1(U_{\mathbb{Q}_p})$  a basis for  $H_{\text{dR}}^1(X_{\mathbb{Q}_p})$ .

Then there exist:

- a matrix  $\Phi \in M_{2g \times 2g}(\mathbb{Q}_p)$ ,
- (overconvergent) functions  $f_1, \dots, f_{2g}$  on some open of  $X_{\mathbb{Q}_p}$ ,

such that

$$F_p^*(\omega_i) = df_i + \sum_{j=1}^{2g} \Phi_{ij} \omega_j \quad \text{for } i = 1, \dots, 2g.$$

We can take  $\omega_1, \dots, \omega_g$  to be a basis for  $H^0(X_{\mathbb{Q}_p}, \Omega^1)$ .

# General integrals

Recall that

$$F_p^*(\omega_i) = df_i + \sum_{j=1}^{2g} \Phi_{ij} \omega_j \quad \text{for } i = 1, \dots, 2g.$$

Assume that  $F_p(P) = P$  and  $F_p(Q) = Q$  (Teichmüller points). No loss of generality, can correct with tiny integrals. Integrating, we find

$$\int_P^Q \omega_i = \int_{F_p(P)}^{F_p(Q)} \omega_i = \int_P^Q F_p^*(\omega_i) = f_i(Q) - f_i(P) + \sum_j \Phi_{ij} \int_P^Q \omega_j.$$

So we can determine the  $\int_P^Q \omega_i$  by solving the linear system

$$(\Phi - I) \int_P^Q \omega_i = f_i(P) - f_i(Q) \quad \text{for } i = 1, \dots, 2g.$$

# Chabauty-Coleman

Assume a point  $b \in X(\mathbb{Q})$  is known and embed  $X \hookrightarrow J$  into its Jacobian by  $P \mapsto (P) - (b)$ .

## Theorem (Chabauty-Coleman)

Let  $r$  denote the Mordell-Weil rank of  $J$  and suppose that  $r < g$ . Then there exists  $\omega \in H^0(X_{\mathbb{Q}_p}, \Omega^1)$  such that  $\int_b^P \omega = 0$  for all  $P \in X(\mathbb{Q})$ .

Sketch of proof.

$$\begin{array}{ccccc} X(\mathbb{Q}) & \longrightarrow & X(\mathbb{Q}_p) & & \\ \downarrow & & \downarrow & \searrow^{AJ_b} & \\ J(\mathbb{Q}) & \longrightarrow & J(\mathbb{Q}_p) & \xrightarrow{D \mapsto \int_D} & H^0(X_{\mathbb{Q}_p}, \Omega^1)^* \end{array}$$

$X(\mathbb{Q})$  lands in a subspace of  $H^0(X_{\mathbb{Q}_p}, \Omega^1)^*$  of dimension at most  $r$ . □



# Implementation

Together with Balakrishnan we have developed and implemented algorithm for computing (single) Coleman integrals and carrying out effective Chabauty on *arbitrary curves*:

[www.github.com/jtuitman/Coleman](https://www.github.com/jtuitman/Coleman)

Note that we need Mordell-Weil rank  $r$  to be known and  $r < g$ .

When  $r \geq g$  there is an extension of the effective Chabauty method by Kim (non-Abelian Chabauty), involving *iterated* Coleman integrals.

Together with Balakrishnan, Dogra, Müller and Vonk we have recently succeeded in applying Kim's method to the split Cartan modular curve of level 13 (also known as the *cursed curve*).

## Example

Let us return to the example  $f(x, y) = y^3 - (x^5 - 2x^4 - 2x^3 - 2x^2 - 3x)$ . The Magma function `RankBounds()` proves that the rank of  $J$  is 1. This uses work of Poonen-Schaefer (1997). Now we use our code:

```
> load "coleman.m";
> Q:=y^3 - (x^5 - 2*x^4 - 2*x^3 - 2*x^2 - 3*x);
> p:=7;
> N:=15;
> data:=coleman_data(Q,p,N);
> Qpoints:=Q_points(data,1000); // PointSearch
> #vanishing_differentials(Qpoints,data:e:=50);
3
> #effective_chabauty(data,1000:e:=50),#Qpoints;
5 5
```

This proves that our list of rational points is complete.